

A Novel Secure Multiparty Algorithms in Horizontally Distributed Database for Fast Distributed Database

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Abstract

We proposed a protocol for protected mining of association rules in horizontally scattered database. The current leading set of rules is that of Kantarcioglu and Clifton. Our procedure, like theirs, is based on the Fast Distributed Mining (FDM) algorithm of Cheung et al. which is an unsecured spread version of the Apriori algorithm. The main ingredient in our procedure are two novel secure multi-party algorithms^[1] — one that computes the union of private subsets that each of the interacting group of actors hold, and another that tests the inclusion of an element held by one actor in a subset held by another. Our protocol offers improved separation with respect to the protocol. In addition, it is simpler and is extensively more efficient in terms of announcement rounds, announcement cost and computational cost.

Keywords: *Privacy Preserving Data Mining; Distributed Computation; Frequent Itemsets; Association Rules.*

I. Introduction

The most precious part of the protocol and its implementation relies upon cryptographic primitives such as commutative encryption, oblivious remove, and confusion functions. The main part of the procedure is a sub-protocol for the secure computation of the union of private subsets that are held by the different group of actors This is also the only part in the procedure in

which the players may extract from their view of the protocol information on other database past what is implied by the final output and their own input. Leakage of information renders the protocol not perfectly protected, the perimeter of the excess information is explicitly bounded It is argued there that such information leakage is safe whence acceptable from a practical point of view. Insufficient security, simplicity and efficiency are not well in the databases, not sure in isolation in an existing system. While our solution is still not

perfectly protected, it leaks excess information only to a little number (three) of probable II. coalitions, unlike the protocol of that discloses information also to some single group of actors. Our protocol may leak is less sensitive than the excess information leaked by the protocol.

II .Problem Statement

The proposed protocol improves upon that in terms of simplicity and efficiency as well as privacy. Our protocol does not depend on commutative encryption and oblivious transfer. We propose here computes a parameterized family of functions, which we call doorstep functions, in which the two great cases communicate to the problems of computing the union and intersection of private subsets.

The excess information that our protocol may leak is less sensitive than the excess information leaked by the protocol This project proposd two novel secure multi-party algorithms — one that computes the union of private subsets that each of the interacting players hold and another that tests the inclusion of an element held by one player in a subset held by another. The problem of secure multiparty computation that we solve here is the set inclusion problem.

Objective : We proposed a protocol for secure mining of association rules in horizontally distributed databases^[2] that improves significantly upon the current leading protocol in terms of privacy^[3] and efficiency. The main ingredient in our proposed protocol is a novel secure multi-party protocol for compute the union (or intersection) of private subsets that each of the interacting players holds.

III. Proposed System

The proposed protocol improves upon that in terms of simplicity and efficiency as well as privacy Our protocol does not depend on commutative encryption and oblivious transfer.

We suggest here computes a parameterized family of function, which we call threshold functions, in which the two extreme cases correspond to the troubles of computing the union and intersection of private subsets. The overload information that our procedure may leak is less sensitive than the excess information leaked by the protocol

The problem of secure multiparty^[4] computation that we solve here is the set inclusion problem

The advantages are as following.

We proposed a procedure for protected mining of society rules in horizontally distributed databases that improves significantly upon the current leading protocol in terms of privacy and efficiency.

The main ingredient in our projected procedure is a novel secure multi-party^[5] protocol for computing the union (or intersection) of private subsets that each of the interacting players holds.

IV. Algorithm as Proposal

Algorithm Analyzed The first pass of the algorithm simply counts item occurrence determine the big 1-itemsets. A successive pass,say pass k, consists of two phase First, the large itemsets L_{k-1} set up in the (k-1)th pass are used to generate the candidate itemsets C_k, using the apriori- production function described in Section 2.1.1. Next, the database is scanned and

the support of candidates in C_k is counted. For fast excluding we need to proficiently determine the candidates in C_k that are contained in a given transaction t . Section 2.1.2 describes the subset function used for this purpose.

Table 1: Notation associated with the candidates.

- 1) $L_1 = \text{flarge 1-itemsets}$;
- 2) for ($k = 2; L_k \neq \emptyset; k++$) do begin
- 3) $C_k = \text{apriori-gen}(L_{k-1})$; // New candidates
- 4) forall transactions $t \in D$ do begin
- 5) $C_t = \text{subset}(C_k, t)$; // Candidates contained in t
- 6) forall candidates $c \in C_t$ do
- 7) $c:\text{count}++$
- 8) end
- 9) $L_k = \{c \in C_k \mid c:\text{count} \geq \text{minsupg}\}$
- 10) end
- 11) Answer = $\bigcup_k L_k$;

V. System Architecture

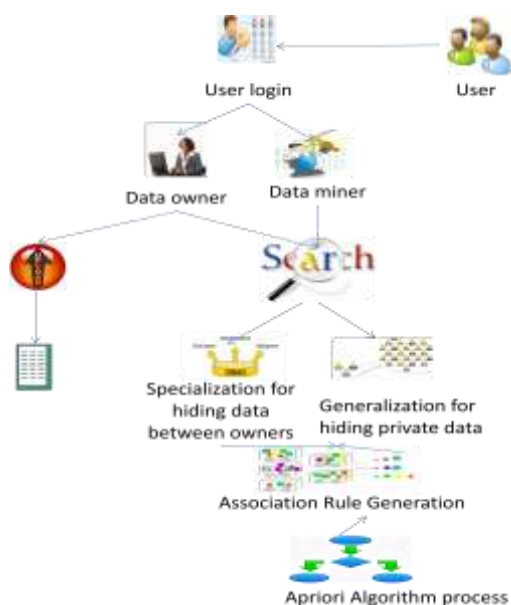


Fig 1: System Architecture

The above system architecture explains the user module which enlists the privacy preserving data mining has considered two connected settings. One, in which the data owner and the data miner are two different individual, and another, in which the data is scattered among several parties who aim to jointly perform data mining on the unified

k-itemset	An itemset having k items.
LK	Set of large k-itemsets (those with minimum support). Each member of this set has two _ elds itemset and ii) support count.
Ck	Set of candidate k-itemsets (potentially large itemsets). Each member of this set has tw: i) itemset and ii) support count.
Ck	Set of candidate k-itemsets when the TI of the generating transactions arekept

corpus of data that they grip In the first location, the goal is to protect the data records from the data miner. Hence, the information holder aims at anonymizing the data prior to its release. The main approach in this framework is to apply data perturbation. He disconcerted data can be used to infer general trends in the data, without informative original documents information. In the second setting, the goal is to perform data mining while protecting^[6] the data records of each of the data owners from the other data owners. The work of the administrator^[8] is to view user details. direction to view the item set based on the user processing details using association role with Apriori algorithm connection rules are if/then

statements that help uncover relationships between seemingly unrelated data in a relational database or other information depository. An example of an association rule would be "If a customer buys a dozen seed, he is 80% likely to also purchase develop. Association set of laws are created by analyzing data for frequent if/then patterns and using the criteria support and confidence to identify the most important associations. Support is an indication of how regularly the items appear in the database. self-confidence indicate the number of period the if/then statement have been found to be true^[7].

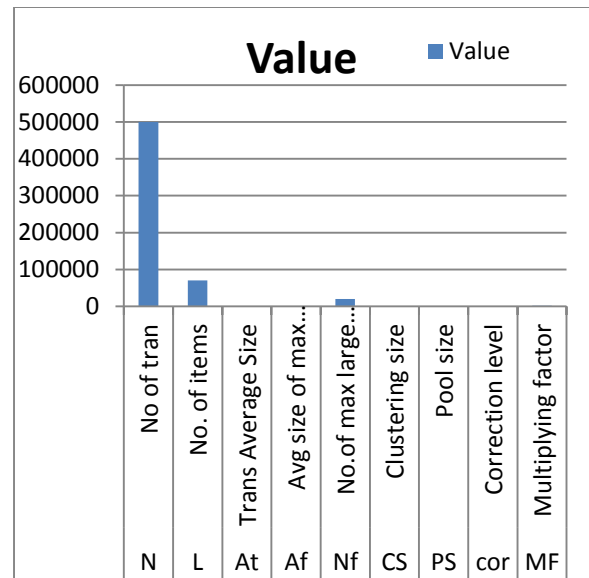


Fig 2. Parameter Value

VI. EXPERIMENTAL RESULTS

Table 1. Parameter For Generating the synthetic Database

Parameter	Generation Method	Value
N	Number of transactions in the whole database	500000
L	Number of items	70000
At	Transaction Average Size	10
Af	Average size of maximum potentially large itemsets	4
Nf	Number of maximum potentially large itemsets	20000
CS	Clustering size	5
PS	Pool size	60
cor	Correction level	0.5
MF	Multiplying factor	1800

We compared the performance of two secure implementations of the FDM algorithm. In the first implementation (denoted FDM-KC), we executed the unification step (Step 4 in FDM) using procedure UNIFI-KC, where the commutative nobody was 1024-bit RSA in the second implementation (denoted FDM) we used our procedure UNIFI, where the keyed-hash purpose was HMAC. In both implementations, we implemented in FDM algorithm in the secure manner that was described. We tested the two implementations with respect to three measures:

- 1) Total working out time of the complete protocols (FDMKC and FDM) over all players. That measure includes the Apriori computation time^{[8][9]}, and the time to identify the globally s -frequent itemsets, as described in Section
- 2) Total working out time of the unification protocols only
3. (The latter two procedures are implemented in the same way in both Protocols FDM-KC and FDM.) (UNIFI-KC and UNIFI) over all players.

4) Total message size. We ran three experiment sets, where each set tested the

Dependence of the above measures on a diverse parameter:

- N — the number of transactions in the unified database,
- M — the number of players, and
- s — the threshold support size.

In our basic configuration, we took $N = 500,000$, $M = 10$ and $s = 0.1$. In the first experiment set, we kept M and s fixed

and tested several values of N . In the second experiment set, we kept N and s fixed and varied M . In the third set, we kept

N and M fixed and varied s . The results in each of those experiment sets are shown above. All experiments were implemented in C# (.net 4) and were executed on an Intel(R) Core(TM)i7-2620M personal computer with a 2.7GHz CPU, 8 GB of RAM, and the 64-bit operating system Windows 7 Professional SP1.

VII. CONCLUSION:

We proposed a protocol for secure mining of association rules in horizontally distributed databases that improves significantly upon the current leading protocol in terms of privacy^{[10][11]} and competence. One of the main ingredients in our proposed protocol is a novel secure multi-party protocol for computing the union (or intersection) of private subsets that each of the interacting group of actors hold. Another ingredient is a procedure that tests the inclusion of an element held by one player in a subset held by another. Those protocols exploit the fact that the underlying problem is of interest only when the number of players is greater than two. One

research problem that this study suggests was described namely, to devise an efficient protocol for inequality verifications that uses the existence of a semihonest third party. Such a protocol might enable to further improve upon the announcement and computational costs of the second and third stages of the protocol, as described earlier. Other research problems that this study suggests is the implementation of the techniques presented here to the problem of distributed association rule mining in the vertical setting, the problem of mining generalized association rules, and the problem of subgroup detection in horizontally partitioned data.

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